

# SOLVING WCSP BY EXTRACTION OF MINIMAL UNSATISFIABLE CORES

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## Introduction

### Context:

- CSP and Weighted CSP frameworks
- WCSP algorithms are often more complex than their CSP counterparts (due to management of costs)

### Goal:

- Benefit from efficient CSP algorithms developed for more than two decades

### Principle:

- Solve WCSP by iteratively generating and solving classical CSPs (greedy approach)
- The sequence of CSPs is enumerated according to an increasing cost order related to the WCSP
- Minimal Unsatisfiable Cores (MUC) identify the soft constraints whose costs must be increased

## Background

### CSP framework:

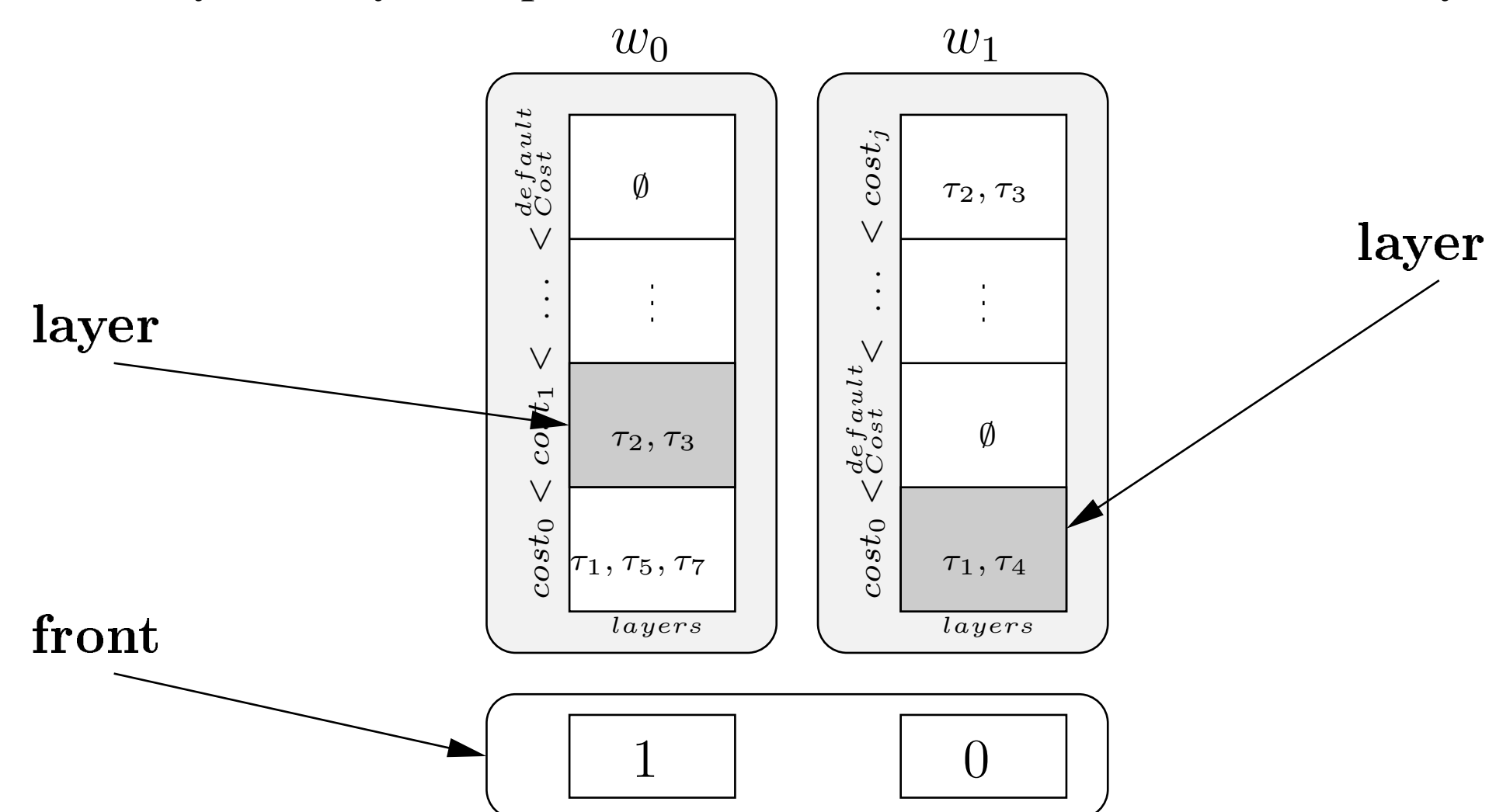
- A CSP is satisfiable iff it admits at least one solution
- An Unsatisfiable Core is an unsatisfiable subset of constraints
- A core is a **Minimal Unsatisfiable Core (MUC)** iff each strict subset is satisfiable

### WCSP framework:

- Extension of CSP
- $(X, C, k)$ :  $C$  is a set of soft constraints (cost functions),  $k > 0$  is either a natural integer or  $+\infty$
- $\forall a, b \in \{0, \dots, k\}, a \oplus b = \min(k, a + b)$
- Goal: find a complete instantiation with minimal cost (optimisation problem)
- The current methods to solve WCSPs: branch and bound tree search combined with the use of soft local consistencies (EDAC, etc. by cost transfer) for estimating minimal costs of sub-problems during search

## Layers and fronts

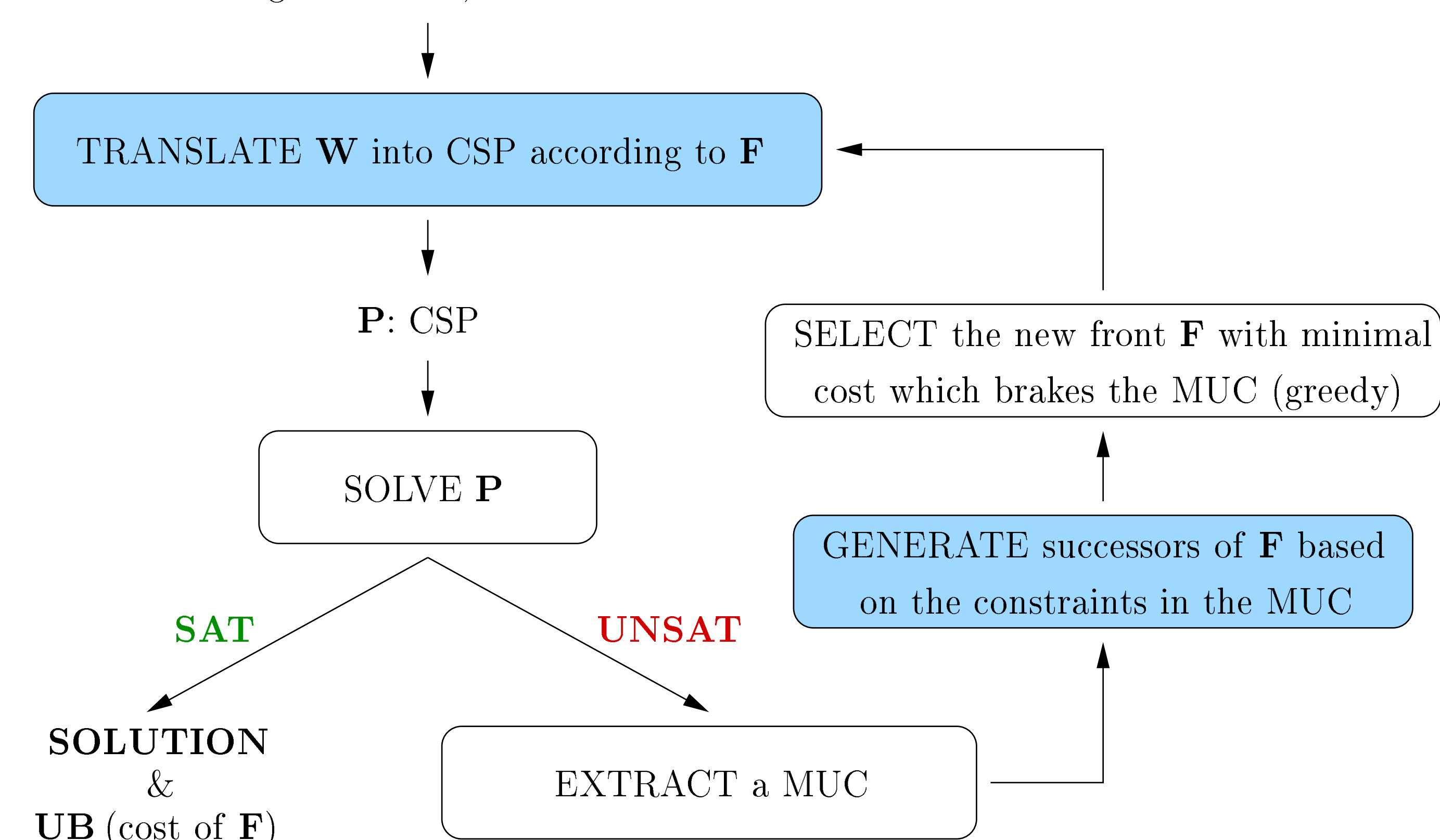
- Focus on soft table constraints (explicit and implicit tuples), but the method can be easily extended to other kinds of constraints
- A **layer** contains all tuples having the same cost
- A **front** (represented by an array  $f$ ) maps each constraint of a WCSP to one of its layers



- **Cost of a front**: sum of costs associated with the selected layers

## General principle of the greedy approach

$W$  : original WCSP,  $F$  : initial front



## Algorithm

GMR ( $W$ : WCSP)

```

foreach  $w \in \text{constraints}(W)$  do
   $f[w] \leftarrow 0$ ;
repeat
   $P \leftarrow \text{toCSP}_{\leq}(W, f)$ ;
   $\text{sol} \leftarrow \text{solveCSP}(P)$ ;
  if  $\text{sol} \neq \perp$  then
    return  $\text{sol}$ ;
  else
     $M \leftarrow \text{extractMUC}(P)$ ;
     $W' \leftarrow \text{restrict}(W, M)$ ;
     $f \leftarrow \text{relax}(W', f)$ ;
until  $\text{sol} \neq \perp$ ;

```

### • $\text{toCSP}_{\leq}(W, f)$

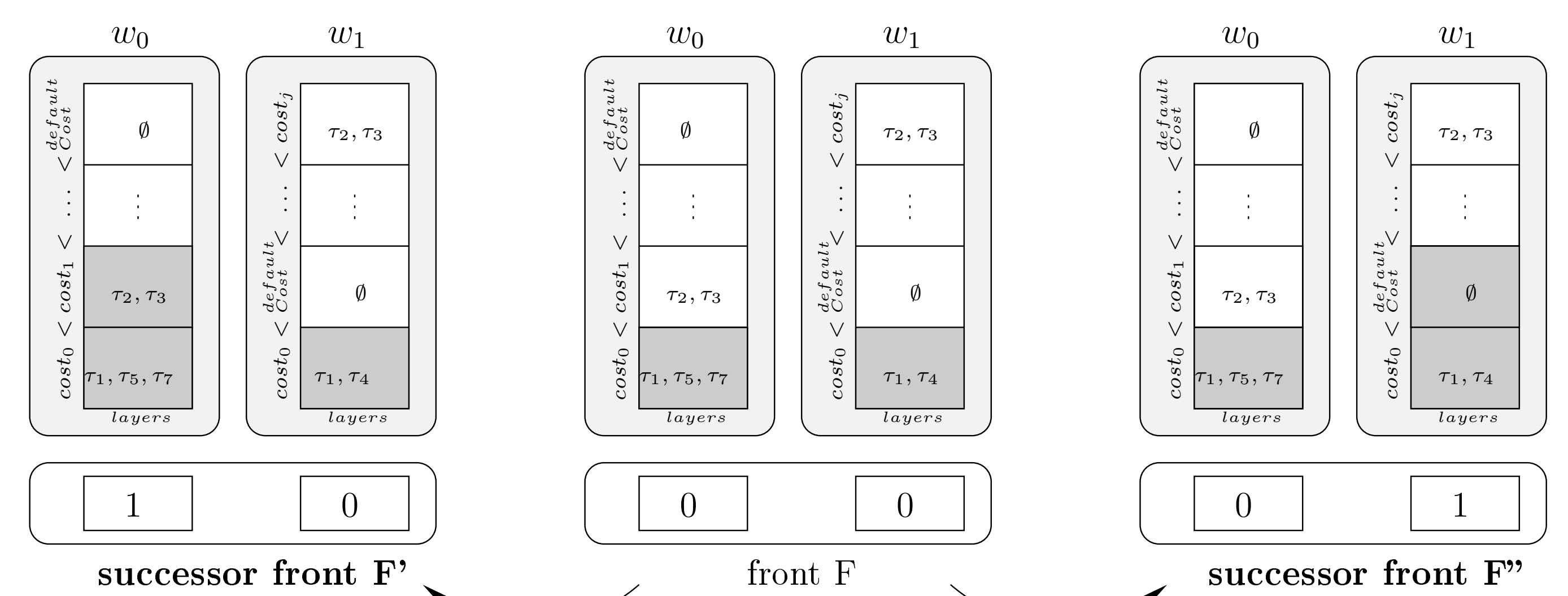
- Translates WCSP into CSP: converting soft constraints into hard constraints according to a front  $f$
- Considering a soft constraint  $w$  and the front  $f$ : a hard constraint is obtained by selecting as allowed tuples in  $w$  the tuples of the layers whose index is **less than or equal to**  $f[w]$
- Representation: extension, positive table (supports) / negative table (conflicts)
- Default cost Layer forbidden (resp. allowed)  $\Rightarrow$  hard positive (resp. negative) constraint

### • $\text{extractMUC}(P)$

- A dichotomic approach is used to extract MUCs of unsatisfiable CSPs

### • $\text{relax}(W', f)$

- MUCs are broken by generating **successors** of the front considering only the constraints of MUCs



- A **successor** of a front differs only by the incrementation of the allowed layers for one constraint

## Experiments

Instances		AbsCon		ToulBar2	
		GMR	EDAC	INCOP	EDAC
spot5/spot5-404	CPU	4.99	⊥	⊥	217
	UB	118	<b>114</b>	<b>114</b>	<b>114</b>
spot5/spot5-412	CPU	18.8	⊥	⊥	⊥
	UB	33,403	43,390	<b>32398</b>	37,399
spot5/spot5-505	CPU	12	⊥	⊥	⊥
	UB	22,266	28,258	<b>21266</b>	25,268
spot5/spot5-509	CPU	32.2	⊥	⊥	⊥
	UB	37469	48,475	<b>37462</b>	46,477
spot5/spot5-1403	CPU	142.5	⊥	⊥	⊥
	UB	<b>481,266</b>	517,260	482267	507,265
celar/graph-05	CPU	16.6	⊥	⊥	0.62
	UB	<b>221</b>	4,645	243	<b>221</b>
celar/scen-06-20	CPU	68.5	⊥	⊥	67.9
	UB	3,402	8,594	3166	<b>3,163</b>
celar/scen-07	CPU	209.9	⊥	⊥	⊥
	UB	426,423	31,230K	<b>394006</b>	505,731

CPU time (in seconds) to prove optimality on various selected instances (time-out of 600 seconds set per instance,  $\perp$ : time-out reached)

## Conclusion and future work

### • Conclusion

- Original greedy approach: solve WCSP through successive resolutions of CSPs
- Focus the cost increase on the sole constraints in the Minimal Unsatisfiable Cores extracted
- Promising results when compared to other state-of-the-art approaches

### • Future work

- Complete approach based on the same principles (work in progress)